

Quantum Mechanics C (Physics 130C) Winter 2015 Worksheet 8

Announcements

- The 130C web site is:

<http://physics.ucsd.edu/~mcgreevy/w15/> .

Please check it regularly! It contains relevant course information!

This week we'll be discussing phase estimation, order finding, and factoring integers.

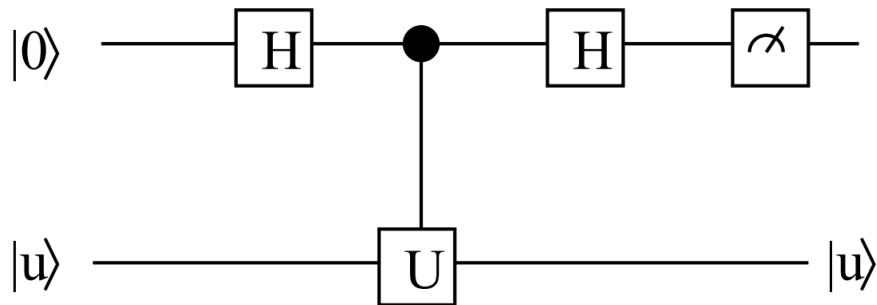
Problems

1. Phase Estimation

Consider a two-qubit system and a unitary operator U with eigenvector $|u\rangle$

By unitarity we know that the eigenvalue has the form $U|u\rangle = e^{i2\pi\theta}|u\rangle$ for some $\theta \in (0, 1)$ which we would like to determine. Challenge: this eigenvalue has norm 1.

So! Consider the following quantum circuit:



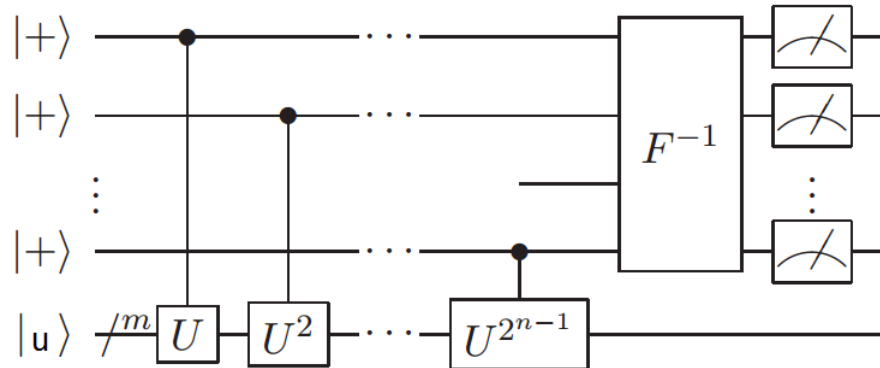
where $\hat{H} = \frac{|0\rangle+|1\rangle}{\sqrt{2}}\langle 0| + \frac{|0\rangle-|1\rangle}{\sqrt{2}}\langle 1|$ is the Hadamard gate and the second line is the Control- U operator $CU = |0\rangle\langle 0| \otimes \mathbb{1} + |1\rangle\langle 1| \otimes U$

- Write down the wavefunction for the qubit pair after each step in the circuit
- After the circuit is applied show that the probability of measuring the first qubit to be $|0\rangle$ is $p = \cos^2(\pi\theta)$. Note the answer if $\theta = \frac{z}{2^T}$ for some integer z .

So by estimating p we can determine θ perfectly in the special case. More generally we could crudely estimate it up to the ambiguity of $\cos^2(\pi\theta) = \cos^2(\pi(1 - \theta))$

Can we do better? Suppose instead of two qubits we have $n + 1$ and an input state $|In\rangle = |00 \cdots 0\rangle|u\rangle$

- (c) Apply the Hadamard gate to each of the first n qubits. What is the result?
 Now consider this improved phase estimation circuit



where F^{-1} is the *quantum* inverse Fourier transformation and the other gates are CU^{2^j} gates applied incrementally to the first n qubits and $j \in \{0, 1, \dots, n-1\}$

- (d) Before the F^{-1} , what is the state of the system? (Hint: What is $CU^{2^j}|+\rangle|u\rangle$?)
 Now we need to tackle this Fourier transformation. For our purposes it is sufficient to define it by

$$F|x\rangle = \frac{1}{\sqrt{N}} \sum_y e^{\frac{2\pi ixy}{N}} |y\rangle \quad (1)$$

- (e) Show that for $\theta = \frac{z}{N}$ for some integer z that the output of the phase estimator circuit gives probability 1 to determine z correctly when measuring the control qubits in the computational basis. What's the probability more generally?

2. Order Finding

Suppose I want to find a factor of an integer N . We'll pick a random integer $0 < a < N$ such that $\text{GCD}[a, N] = 1$ and determine the integer r such that $a^r = 1 \pmod{N}$

Without proof, I claim that with $> \frac{1}{2}$ probability r is even and $a^{\frac{r}{2}} \not\equiv \pm 1 \pmod{N}$

Given this N divides $a^r - 1$ which implies N divides $(a^{\frac{r}{2}} + 1)(a^{\frac{r}{2}} - 1)$ but not each independently. Thus atleast one of $(a^{\frac{r}{2}} \pm 1)$ contains a factor of N which we pick out by computing $\text{GCD}[a^{\frac{r}{2}} + 1, N]$

- (a) Consider the case of $N = 15$ and $a = 4$.

Hopefully the above convinces you that order finding is the thing to do. To accomplish this consider the following unitary operator:

$$U_a|x\rangle = |ax \pmod{N}\rangle \quad 0 \leq x \leq N \quad (2)$$

- (b) Show that $\lambda_k = e^{\frac{2\pi ik}{r}} \equiv \omega^k$ are eigenvalues of U_a for $0 \leq k < r$

- (c) Show that $|\phi_k\rangle = \frac{1}{\sqrt{r}} \sum_{\ell=0}^{r-1} e^{-\frac{2\pi ik\ell}{r}} |a^\ell \pmod{N}\rangle$ are eigenvectors of U_a

But now suppose you were handed $|\phi_k\rangle$ delicately prepared. We could use the phase estimation circuit above to estimate $\theta = \frac{k}{r} \approx \frac{z}{2^n}$ which, as long as k and r are not relatively prime¹, we get r !

Now it is a justified complaint that, without knowing r , constructing $|\phi_k\rangle$ would be impossible.

However, superpositions of these eigenvectors aren't hard to make.

(d) Show $|1\rangle = \frac{1}{\sqrt{r}} \sum_k |\phi_k\rangle$

So using $|1\rangle$ as our input our measurement still produces, with good probability, a rational approximation to $\frac{k}{r}$ but for some uniformly random k . Is it good enough?

As it turns out yes! But that's beyond today.

¹For large r this get's unlikely, another claim without proof.