# Gapped domain walls, symmetry-protected topological phases, and fault-tolerant logical gates in topological color code

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We discuss a connection between the gapped domain walls, symmetry-protected topological phases, and fault-tolerant logical gates in topological color codes based on [Phys. Rev. B 91, 245131]. It was found that applying a *d*-dim transversal operator leads to d - 1-dim excitations characterized by bosonic symmetry-protected topological (SPT) wave functions, and these SPT-excitations can be realized as transparent gapped domain walls in the color code. The connection may be generalized to a large class of topological quantum codes and topological quantum field theories.

### I. INTRODUCTION

Symmetry-protected topological (SPT) quantum phases are described by short-range entangled states that cannot be smoothly connected to trivial product states in the presence of symmetries [1-3]. The study of SPT phases has enriched quantum many-body physics in many aspects, yet their implication in quantum information science is less explored. In Ref. [4, 5], Yoshida explores an intriguing connection between SPT phases and fault-tolerant logical gates in topological quantum codes. Such a connection is notable since one may apply the known classification of SPT phases to classify or construct fault-tolerant logical gates, which is essential in topological quantum computation. Furthermore, Yoshida also pointed out that gapped boundaries and domain walls[6], another important subject in the study of topological order, can be constructed given the knowledge of fault-tolerant logical gates. As such, a connection between SPT phases, fault-tolerant logical gates, and gapped boundaries can be established.

In this report, we will provide a short review regarding the aforementioned connection. We will mainly follow Ref.[4], which focuses on d-dim topological color codes[7, 8], and the generalization to the quantum double model discussed in Ref.[5] will not be discussed here.

### II. TOPOLOGICAL COLOR CODE

A two-dimensional color code can be defined on any threevalent and three-colorable lattice, where each vertex accommodates a qubit. One common choice is the hexagonal lattice with the Hamiltonian defined as

$$H = -\sum_{P} S_{P}^{(X)} - \sum_{P} S_{P}^{(Z)}.$$
 (1)

*P* labels a plaquette, and  $S_P^{(X)}$ ,  $S_P^{(Z)}$  are products of Pauli-X, Z operators acting on all qubits on the plaquette *P* (see Fig.1). This is a stabilizer Hamiltonian since every term commutes with each other, and correspondingly, a ground state  $|\psi\rangle$  satisfies  $S_P^{(X)} |\psi\rangle = S_P^{(Z)} |\psi\rangle = |\psi\rangle$ . Being a topological code, *H* supports anyonic excitations, which live on two ends of string operators. To construct these excitations, we now assign colors *A*, *B*, and *C* such that two plaquettes that share an edge have different colors. It follows that an edge can also have the



FIG. 1: Left: local terms  $S_P^{(Z)}$  and  $S_P^{(X)}$  in the color code Hamiltonian (Eq.1). Right: a string operator  $\gamma^{AB}$  creates a pair of excitations on its two ends (the shaded plaquettes).

color *AB*, *BC*, or *CA* that is determined by its two neighboring plaquettes. For a set of edges of color *AB* which form a one-dimensional line  $\gamma^{AB}$  (see Fig.1), one can define

$$\overline{X^{AB}}|_{\gamma^{AB}} = \prod_{j \in \gamma^{AB}} X_j, \quad \overline{Z^{AB}}|_{\gamma^{AB}} = \prod_{j \in \gamma^{AB}} Z_j, \qquad (2)$$

and it is not hard to see that such string operators commute with all the terms in *H*, except for the two plaquette terms on their boundary. Correspondingly, applying  $\overline{X^{AB}}|_{\gamma^{AB}}$  and  $\overline{Z^{AB}}|_{\gamma^{AB}}$  on a ground state  $|\psi\rangle$  creates the magnetic fluxes  $m_C$ and electric charges  $e_C$  respectively, which can be expressed as

$$\overline{X^{AB}}|_{\gamma^{AB}} \to m_C, \quad \overline{Z^{AB}}|_{\gamma^{AB}} \to e_C. \tag{3}$$

Similarly, one can construct other types of string operators that create anyons:  $\overline{X^{BC}}|_{\gamma^{BC}} \to m_A$ ,  $\overline{Z^{BC}}|_{\gamma^{BC}} \to e_A$ ,  $\overline{X^{CA}}|_{\gamma^{CA}} \to m_B$ , and  $\overline{Z^{CA}}|_{\gamma^{CA}} \to e_B$ . Importantly, the excitations are not independent from each other since applying a single Pauli-X/Z creates the composite excitations  $m_A m_B m_C / e_A e_B e_C$ , implying the existence of the fusion channel  $m_A \times m_B \times m_C = 1$  and  $e_A \times e_B \times e_C = 1$ . In addition to the string operators, there also exist membrane-like operators as follows. Considering the Hadamard operator  $\mathcal{H}$ , which exchanges a Pauli-X and a Pauli-Z, i.e.  $\mathcal{H}X\mathcal{H}^{\dagger} = Z$  and  $\mathcal{H}Z\mathcal{H}^{\dagger} = X$ , one can define the membrane operator

$$\overline{\mathcal{H}} = \prod_{j} \mathcal{H}_{j},\tag{4}$$

which commutes with H but has the non-trivial operation for exchanging magnetic fluxes and electric charges. To construct another non-trivial membrane operator, one first divides the (bipartite) hexagonal lattice into two sublattices T and  $T^c$  so that the nearest neighbors of every site in T belong to  $T^c$  and vice versa, and then a membrane operator  $\overline{R}$  can be defined as

$$\overline{R} = \prod_{j \in T} R_j \prod_{i \in T^c} (R_j)^{-1}, \qquad (5)$$

where  $R = \sqrt{Z} = \text{diag}(1, i)$  is a phase gate which exchanges Pauli-X and Y as  $RXR^{\dagger} = Y$  and  $RYR^{\dagger} = -X$ . Crucially, although  $\overline{R}$  does not commute with the Hamiltonian *H*, it commutes with the projector to the ground state subspace of *H*, and can implement the nontrivial operation on anyons:

$$e_A \to e_A, e_B \to e_B, m_A \to m_A e_A, m_B \to m_B e_B.$$
 (6)

Having introduced the basics of the color code Hamiltonian, we now discuss how SPT orders naturally arise in the application of membrane operators.

### III. SPT ORDER FROM FAULT-TOLERANT LOGICAL GATES

Here we will show that restricting a membrane operator in a subregion *V* induces a loop-like excitation (on its boundary  $\partial V$ ) characterized by a non-trivial SPT order. First, let's consider the flux-free subspace  $\mathcal{H}_{no-flux}$  where  $|\psi\rangle \in \mathcal{H}_{no-flux}$  satisfies  $S_P^{(Z)} |\psi\rangle = |\psi\rangle$ . One can define an excitation basis to encode the location of excitations in  $S_P^{(X)}$ :

$$S_{P_j}^{(X)} \left| \tilde{p}_1, \cdots, \tilde{p}_{n_0} \right\rangle = (1 - 2p_j) \left| \tilde{p}_1, \cdots, \tilde{p}_{n_0} \right\rangle, \tag{7}$$

where  $\tilde{p}_j$  can be 0 or 1, corresponding to the absence or presence of the excitation on the plaquette  $\tilde{p}_j$ , and  $n_0$  denotes the total number of plaquettes on the lattice. Since a ground state  $|\psi_{gs}\rangle$  is invariant under the application of the membrane operators  $\overline{R}$ , restricting  $\overline{R}$  on a subregion V creates a loop-like excitation on the boundary of V (see Fig.2), and the corresponding wave function  $|\psi_V\rangle = \overline{R}|_V |\psi_{gs}\rangle$  can be conveniently expressed in the excitation basis as

$$|\psi_V\rangle \to |\psi_{\partial V}\rangle \otimes \left|\tilde{0}, \cdots, \tilde{0}\right\rangle,$$
 (8)

where  $|\psi_{\partial V}\rangle$  denotes the boundary excitations on the plaquettes  $A_1, B_1, \dots, A_n, B_n$ :

$$|\psi_{\partial V}\rangle = \sum_{\tilde{p}_{A_1}, \tilde{p}_{B_1}, \cdots \tilde{p}_{A_n}, \tilde{p}_{B_n}} \lambda(\{\tilde{p}_j\}) \left| \tilde{p}_{A_1}, \tilde{p}_{B_1}, \cdots \tilde{p}_{A_n}, \tilde{p}_{B_n} \right\rangle$$
(9)

and  $|\tilde{0}, \dots, \tilde{0}\rangle$  denotes the rest (un-excited) plaquettes. Notably, the state  $|\psi_{\partial V}\rangle$  is the exact ground state wave function of a 1d cluster state Hamiltonian, which exhibits a non-trivial  $Z_2 \times Z_2$  SPT order. While it is not hard to see that the  $Z_2 \times Z_2$ symmetry arises simply from the number parity conservation



FIG. 2: A loop-like excitation with the  $Z_2 \times Z_2$  SPT order (Eq.9) supported on the boundary of a membrane operator. The phase gate operators  $R(R^{-1})$  are applied on filled circles (filled double circles).

of the electric charges  $e_A$  and  $e_B$ :  $N_A = N_B = 0 \mod 2$ , it remains non-trivial that the boundary excitation corresponds to an SPT phase. The central idea is that the preparation of the boundary loop excitation cannot be obtained by just applying a local unitary transformation localized on the boundary. To see this, one can imagine transporting a magnetic flux *m* from the region outside of *V* to inside of *V*. Based on Eq.6, the *m* flux will be transformed into a pair of *e* charge and *m* flux, and this implies one must apply a unitary operation on all qubits in *V* to realize the boundary excitation, hence giving rise to the non-trivial SPT order. In particular, the application of the *R* phase gate is essentially a symmetry-protected quantum circuit for preparing the SPT wave functions.

## IV. GAPPED DOMAIN WALLS FROM FAULT-TOLERANT LOGICAL GATES

We have seen that applying a membrane operator restricted in a subregion V for a ground state creates fluctuating charges with a non-trivial SPT order on the boundary of V. Now we show that applying such operators to transform Hamiltonian allows to create gapped domain walls, which establishes a connection between the classification of gapped domain walls and fault-tolerant logical gates.

To start, one can split the lattice into the left part and the right part, and perform a transformation on the color code Hamiltonian using the Hadamard gates restricted in the right of the lattice (i.e.  $\mathcal{H}|_R = \prod_{j \in R} \mathcal{H}_j$ ). It follows that the transformed Hamiltonian  $\tilde{H} = \overline{\mathcal{H}}|_R H(\overline{\mathcal{H}}|_R)^{\dagger}$  reads  $\tilde{H} = H_L + H_R + H_{LR}$ , where  $\tilde{H}$  differ from H only in the terms localized on the boundary between L and R. Crucially, since  $\mathcal{H}|_R$  implements a unitary transformation,  $\tilde{H}$  remains gapped, and  $H_{LR}$  can be regarded as a gapped domain wall, across which anyons exchange *m*-fluxes and *e*-charges:  $(e_A|m_A), (m_A|e_A), (e_B|m_B), (m_B|e_B).$ Note that such a domain wall is transparent in the sense that a single anyon cannot condense on the domain wall. One can follow a similar strategy to construct another gapped domain wall using the phase gate R, which induces the following anyon exchange rule  $(m_A|e_Am_A)$ ,  $(e_A|e_A)$ ,  $(m_B|e_Bm_B)$ ,  $(e_B, m_B)$  (see Fig.3). More broadly, one can show that the membrane opera-



FIG. 3: A transparent gapped domain wall obtained by restricting the membrane operator  $\overline{R}$  in the right.

tors associated with non-trivial automorphisms among anyon labels always lead to transparent gapped domain walls in (2 + 1)-dimensional TQFTs.

#### V. SUMMARY AND DISCUSSION

We have presented a short review regarding the connection between SPT phases, gapped domain walls, fault-tolerant logical gates in the two-dimensional topological color code. Such an observation in fact can be generalized to the *d*-dim color code[4]. In particular, applying  $R_d$  phase gates  $(R_d = \text{diag}(1, \exp(i\pi/2^{d-1})))$  on all qubits in a connected subregion *V* creates d - 1 dim excitations characterized by a bosonic SPT order with  $(Z_2)^{\otimes d}$  symmetry, and this SPT order in turn characterizes the gapped domain walls in the *d*-dim color codes.

Finally, we note that the aforementioned connection has been generalized to d-dim quantum double models[5], where it was found that using d-cocycle functions, one can construct the gapped boundaries/domain walls, fault-tolerantly implementable logical gates, and excitations characterized by an SPT order.

Yoshida's works [4, 5] have motivated some interesting questions. A natural future direction is generalizing the discussion of global on-site symmetries of SPT orders to q-form symmetries, where symmetry operators are codimension-1 objects, and exploring the implication to gapped boundaries and faulttolerant logical gates. Similarly, one may employ SPT orders with fractal-like symmetries to explore novel types of gapped boundaries and logical gates.

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